

Location and domination in graphs.

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Outline

Detection devices and graphs

Parameters

Location, Domination and Location-domination
Identifying codes and watching systems

Properties

Values
Bounds
Extremal values
Realization Theorems

Location-Domination in G and \overline{G}

General results
Block-cactus
Bipartite graphs
Global Location-domination

Detection devices and graphs

Graphs are used to model some problems:

detection devices located at some vertices
to detect/locate an intruder in some vertex,
...of course with a small number of detectors

- Detection: is there any intruder?
- Location: where is the intruder?

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- Detection: is there any intruder?

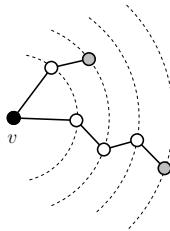
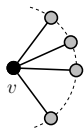
Dominating sets

- Location: where is the intruder?

Locating sets

Restrictions on detection devices

- Detect if there is an intruder in its neighborhood $\rightarrow 0, 1$
- Detect if there is an intruder at distance $\leq k \rightarrow 0, 1$
- Detect if there is an intruder at distance $= k \rightarrow k$



Restrictions on detection devices

- At most one detection device at a vertex
- One or more detection devices at a vertex
- Detect an intruder located at any vertex of the graph
- Detect an intruder in a subset of vertices

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Questions

- Bounds relating different parameters
- Bounds relating order, size, diameter, Δ , δ ,...
- Values on some families: complete graphs, paths, cycles, wheels, bipartite graphs, trees,...
- Extremal values
- Realization type results
- Graph operations: Cartesian product, strong product, complement,...
- Nordhaus-Gaddum type bounds: $p(G) + p(\overline{G})$

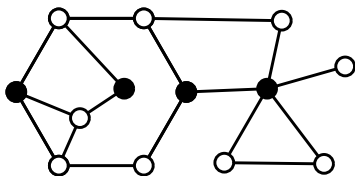
Graphs

$G = (V, E)$ graph,

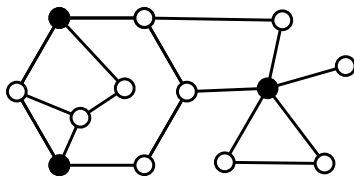
- \overline{G} , *complement of G*
- $N(v) = \{u : uv \in E\}$, *open neighborhood*
- $N[v] = \{v\} \cup N(v)$, *closed neighborhood*
- *true-twin vertices*: $u, v \in V$ such that $N[u] = N[v]$
- *false-twin vertices*: $u, v \in V$ such that $N(u) = N(v)$

Domination

- *Dominating set* of G , $S \subseteq V$:
for all $v \in V \setminus S$, $S \cap N(v) \neq \emptyset$
- *Domination number* of G , $\gamma(G)$:
minimum size of a dominating set of G



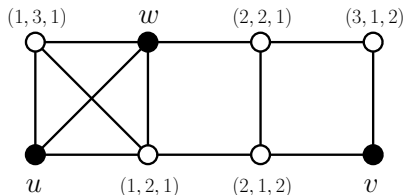
Dominating set



$\gamma(G) = 3$

Location [Slater (1975), Harary and Melter (1976)]

- *Locating set/Resolving set* of G , $S \subseteq V$:
every vertex is uniquely determined by its vector of distances to the vertices of S
- *Location number/Metric dimension* of G , $\beta(G)$:
minimum size of a locating set of G
- *Locating code/Metric basis* of G :
locating and dominating set of minimum size



Location and domination [Henning and Oellermann, 2004]

- *Locating and dominating set (MLD-set) of G , $S \subseteq V$:*
 - dominating set of G
 - locating set of G
- *Location and domination number, $\eta(G)$:*
minimum size of a locating and dominating set of G
- *Locating and dominating code of G :*
locating and dominating set of minimum size

$$\max\{\gamma(G), \beta(G)\} \leq \eta(G) \leq \gamma(G) + \beta(G)$$

Location-domination [Slater, 1988]

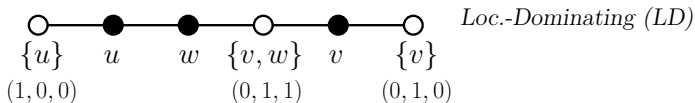
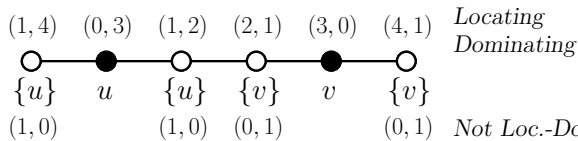
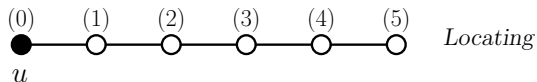
- *Locating-dominating set (LD-set)* of G , $S \subseteq V$:
 - dominating set of G
 - $N_G(u) \cap S \neq N_G(v) \cap S$, if $u, v \in V \setminus S$, $u \neq v$
- *Location-domination number*, $\lambda(G)$:
minimum size of a locating-dominating set of G
- *Locating-dominating code (LD-code)* of G :
LD-set of minimum size

Location-Domination

- ▶ S LD-set $\Rightarrow S$ dominating set
- ▶ S LD-set $\Rightarrow S$ locating set
- ▶ $\gamma(G) \leq \lambda(G)$
- ▶ $\beta(G) \leq \lambda(G)$

$$\max\{\gamma(G), \beta(G)\} \leq \eta(G) \leq \lambda(G)$$

Location and Domination/Location-Domination



Location and Domination/Location-Domination

$$S = \{u_1, \dots, u_r\}$$

$$x \in V \setminus S \longrightarrow \ell(x) = (x_1, \dots, x_r)$$

Param.	x_i	Conditions
γ	$\begin{cases} 1, & \text{if } x \in N(u_i); \\ 0, & \text{otherwise.} \end{cases}$	$\exists x_i = 1$
β	$d(x, u_i)$	$\ell(x) \neq \ell(y), \text{ if } x \neq y$
η	$d(x, u_i)$	$\exists x_i = 1 \quad \ell(x) \neq \ell(y), \text{ if } x \neq y$
λ	$\begin{cases} 1, & \text{if } x \in N(u_i); \\ 0, & \text{otherwise.} \end{cases}$	$\exists x_i = 1 \quad \ell(x) \neq \ell(y), \text{ if } x \neq y$

Identifying codes [Karpovsky, Chakrabarty and Levitin, 1998]

- *Identifying set* of G , $S \subseteq V$:
 - dominating set of G
 - $N_G[u] \cap S \neq N_G[v] \cap S$, if $u, v \in V$, $u \neq v$
- *Identifying number*, $\iota(G)$:
minimum size of an identifying set of G
- *Identifying code* of G :
identifying set of minimum size

Identifying codes exist only in true-twin free graphs

$$\max\{\gamma(G), \beta(G)\} \leq \eta(G) \leq \lambda(G) \leq \iota(G)$$

Watching systems [Auger, Charon, Hudry and Lobstein, 2013]

- *watcher*: $w_i = (v_i, A_i)$, where $w_i = (v_i, A_i)$, $v_i \in V$, $A_i \subseteq N[v_i]$
- *label* of $u \in V$: $L(u) = \{w_i : u \in A_i\}$
- *watching system*: $S = \{w_i : i \in I\}$ such that
 - $L(u) \neq \emptyset$ for all $u \in V$
 - $L(u) \neq L(v)$ if $u \neq v$
- *watching number*, $\omega(G)$:
minimum size of a watching system set of G

Watching systems exist in all graphs

$$\omega(G) \leq \gamma(G) \lceil \log_2(\Delta + 2) \rceil$$

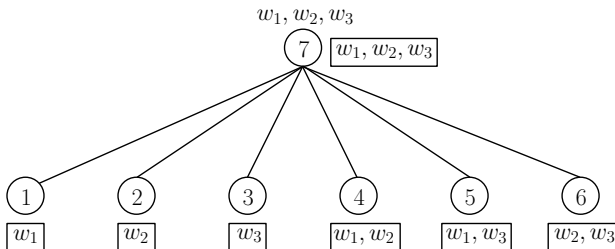
$$\omega(G) \leq \iota(G) \text{ if } G \text{ is twin-free}$$

Watching number: example

$$G = K_{1,6}: i(G) = 6, \omega(G) = 3$$

$$W = \{w_1, w_2, w_3\}, I(w_i) = 7$$

$$A(w_1) = \{1, 4, 5, 7\}, A(w_2) = \{2, 4, 6, 7\}, A(w_3) = \{3, 5, 6, 7\}$$



Some families

	P_n ($n \geq 4$)	C_n ($n \geq 3$)	K_n ($n \geq 2$)	$K_{1,n-1}$ ($n \geq 3$)	$K_{r,s}$ ($2 \leq r \leq s$)
$\gamma(G)$	$\lceil \frac{n}{3} \rceil$	$\lceil \frac{n}{3} \rceil$	1	1	2
$\beta(G)$	1	2	$n - 1$	$n - 2$	$n - 2$
$\eta(G)$	$\lceil \frac{n}{3} \rceil$	$\lceil \frac{n}{3} \rceil$	$n - 1$	$n - 1$	$n - 2$
$\lambda(G)$	$\lceil \frac{2n}{5} \rceil$	$\lceil \frac{2n}{5} \rceil$	$n - 1$	$n - 1$	$n - 2$
	P_n ($n \geq 4$)	C_n ($n \geq 7$)	K_n ($n \geq 2$)	$K_{1,n-1}$ ($n \geq 3$)	$K_{r,s}$ ($2 \leq r \leq s$)
$\iota(G)$	$\lceil \frac{n+1}{2} \rceil$	$3 \lceil \frac{n}{2} \rceil - n$	—	$n - 1$	$n - 2$
$\omega(G)$	$\lceil \frac{n+1}{2} \rceil$	$\lceil \frac{n}{2} \rceil$	$\lceil \log_2(n+1) \rceil$	$\lceil \log_2(n+1) \rceil$	(*)

(*) [Hernando, Mora and Pelayo, 2012]

Bounds

$$G = (V, E), |V| = n, \text{diam}(G) = D \geq 3,$$

- $\beta + D \leq n \leq \left(\left\lfloor \frac{2D}{3} \right\rfloor + 1 \right)^\beta + \beta \sum_{i=1}^{\lceil D/3 \rceil} (2i-1)^{\beta-1}$

[Hernando, Mora, Pelayo, Seara and Wood, 2010]

- $\eta + \left\lfloor \frac{2D}{3} \right\rfloor \leq n \leq \eta(3^{\eta-1} + 1)$, if $G \not\cong K_{1, n-1}$

[Cáceres, Hernando, Mora, Pelayo and Puertas, 2013]

- $\lambda + \left\lfloor \frac{3D-1}{5} \right\rfloor \leq n \leq \lambda + 2^\lambda - 1$

[Cáceres, Hernando, Mora, Pelayo and Puertas, 2013]

All bounds are tight!

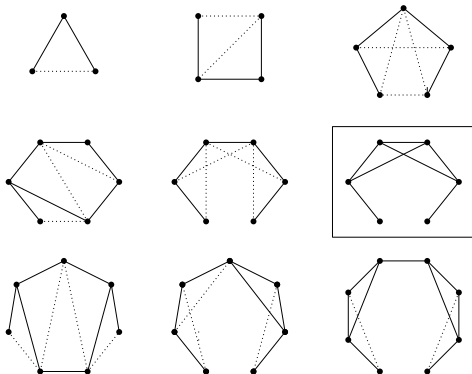
Bounds

$$G = (V, E), |V| = n$$

- $\iota + 1 \leq n \leq 2^\iota - 1$, if G is true-twin free
[Karpovsky, Chakrabarty and Levitin, 1998]
- $n \leq 2^\omega - 1$
[Auger, Charon, Hudry and Lobstein, 2010]
- $\omega(G) \leq \frac{2n}{3}$, if G is a connected graph of order 3 or ≥ 5
[Auger, Charon, Hudry and Lobstein, 2010]

Small values of $\eta(G)$ and $\lambda(G)$

- There are 51 non isomorphic graphs satisfying $\eta(G) = 2$
 - There are 16 non isomorphic graphs satisfying $\lambda(G) = 2$
- [Cáceres, Hernando, Mora, Pelayo and Puertas, 2013]



Large values of $\eta(G)$ and $\lambda(G)$

- $\eta(G) = n - 1 \Leftrightarrow G \cong K_n, K_{1,n-1}$
[Henning and Oellermann, 2004]
- $\lambda(G) = n - 1 \Leftrightarrow G \cong K_n, K_{1,n-1}$
[Slater, 1988]
- Graphs satisfying $\eta(G) = n - 2$
[Henning and Oellermann, 2004]
- Graphs satisfying $\lambda(G) = n - 2$:
 $\eta(G) = n - 2 \iff \lambda(G) = n - 2$
[Cáceres, Hernando, Mora, Pelayo and Puertas, 2013]

Realization Theorems

[Cáceres, Hernando, Mora, Pelayo and Puertas, 2013]

- G graph, $\max\{\gamma(G), \beta(G)\} \leq \eta(G) \leq \gamma(G) + \beta(G)$

$a, b, c \in \mathbb{N}$, $\max\{a, b\} \leq c \leq a + b \Rightarrow$
 $\exists G$ graph satisfying $\gamma(G) = a, \beta(G) = b, \eta(G) = c$,
 except if $1 = b < a < c = a + 1$

- T tree, $|V(T)| \geq 3$, $T \not\cong P_6$ $\eta(T) \leq \lambda(T) \leq 2\eta(T) - 2$

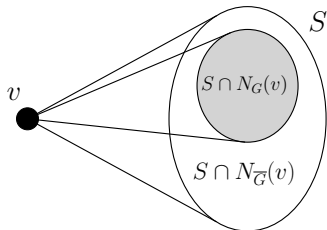
$a, b \in \mathbb{N}$, $3 \leq a \leq b \leq 2a - 2 \Rightarrow$
 $\exists T$ tree satisfying $\eta(T) = a, \lambda(T) = b$

$\lambda(G)$ versus $\lambda(\overline{G})$ [Hernando, Mora and Pelayo (2013)]

S LD-set of $G = (V, E)$:

- dominating set of G
- $N_G(u) \cap S \neq N_G(v) \cap S$, if $u, v \in V \setminus S, u \neq v$

► $N_G(x) \cap S \neq N_G(y) \cap S \Leftrightarrow N_{\overline{G}}(x) \cap S \neq N_{\overline{G}}(y) \cap S$



$\lambda(G)$ versus $\lambda(\overline{G})$

S LD-set of G , then

- ▶ S LD-set of $\overline{G} \Leftrightarrow S$ dominating set of \overline{G} .
- ▶ S LD-set of $\overline{G} \Leftrightarrow \nexists u \in V \setminus S, u$ dominates S in G .
- ▶ $\exists u \in V \setminus S, u$ dominates S in G
 $\Rightarrow S \cup \{u\}$ is an LD-set of \overline{G}

$$|\lambda(G) - \lambda(\overline{G})| \leq 1$$

Global LD-sets

- *global LD-set of G* : LD-set of G and of \overline{G}
i.e., $\nexists u \in V \setminus S, u$ dominates S in G .
- ▶ G contains a global LD-code $\Rightarrow \lambda(\overline{G}) \leq \lambda(G)$

Graphs satisfying $\lambda(\overline{G}) = \lambda(G) + 1$?

Global LD-sets

G contains a non-global LD-set $S \Rightarrow$

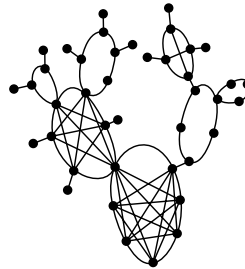
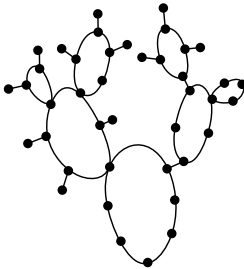
- ▶ $\exists u \in V \setminus S$, u dominates S
- ▶ $\text{ecc}(u) \leq 2$, $\text{rad}(G) \leq 2$, $\text{diam}(G) \leq 4$

$\lambda(\overline{G}) = \lambda(G) + 1 \Rightarrow$

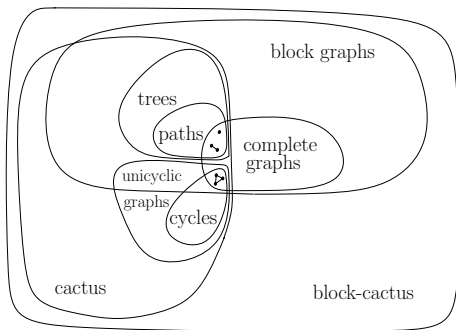
- \Rightarrow every LD-code of G is non-global
- $\Rightarrow G$ connected and $\text{diam}(G) \leq 4$

Block-cactus

- Block of $G = (V, E)$: maximally connected subgraph with no cut vertices
- Cactus: connected graph s.t. all blocks are cycles or K_2 , i.e., there is no edge lying on two different cycles
- Block-cactus: connected graph s.t. all blocks are cycles or complete graphs



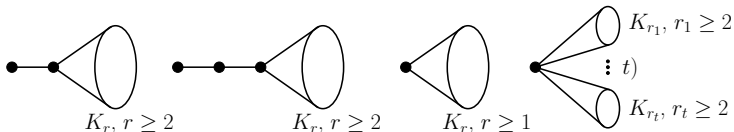
Families of block-cactus



Block-cactus

Block-cactus s.t. $\lambda(\overline{G}) = \lambda(G) + 1$

G is one of the following graphs:



► T tree of order $\geq 3 \Rightarrow \lambda(\overline{T}) \leq \lambda(T)$

Bipartite graphs

Bipartite graphs s.t. $\lambda(\overline{G}) = \lambda(G) + 1$

$G = (V, E)$ $V = V_1 \cup V_2$, $|V_1| = r$, $|V_2| = s$, $2 \leq r \leq s$

▶ $r = 1, 2 \Rightarrow \lambda(\overline{G}) \leq \lambda(G)$

▶ S LD-code of G

- $S \cap V_1 \neq \emptyset$ and $S \cap V_2 \neq \emptyset \Rightarrow \lambda(\overline{G}) \leq \lambda(G)$.
- $S = V_2$ and $r < s \Rightarrow \lambda(\overline{G}) \leq \lambda(G)$.

Bipartite graphs satisfying $\lambda(\overline{G}) = \lambda(G) + 1$

$$G = (V, E), V = V_1 \cup V_2, |V_1| = r, |V_2| = s, 3 \leq r \leq s$$

Theorem

$$\lambda(\overline{G}) = \lambda(G) + 1 \Rightarrow \frac{3r}{2} \leq s \leq 2^r - 1$$

Proof.

$$\lambda(\overline{G}) = \lambda(G) + 1 \Rightarrow V_1 \text{ LD-code of } G$$

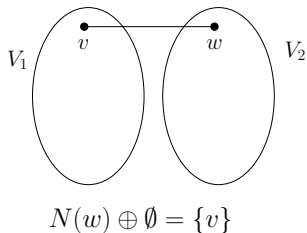
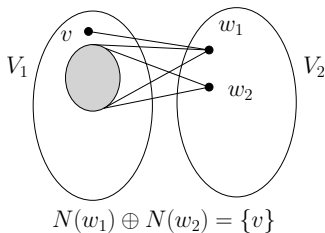
$$s \leq 2^r - 1$$

$$V_1 \text{ LD-code} \Rightarrow s \leq 2^r - 1$$

Proof of $\frac{3r}{2} \leq s$

$$\frac{3r}{2} \leq s$$

$\forall v \in V_1, V_1 \setminus \{v\}$ is not an LD-set:



Proof of $\frac{3r}{2} \leq s$: the graph G^*

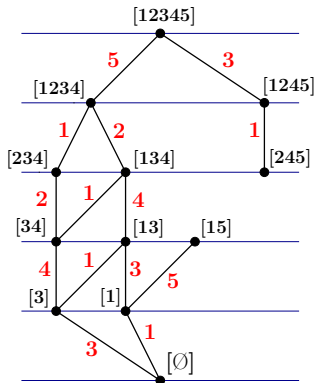
G^* edge-labeled graph associated to G :

- $V(G^*) = V_2 \cup \{w_0\}$, $w_0 \notin V_2$ and define $N(w_0) = \emptyset$
- $w_h w_k \in E(G^*) \Leftrightarrow N(w_h) \oplus N(w_k) = \{v\}$ for some $v \in V_1$
- $\ell(w_h w_k) = N(w_h) \oplus N(w_k) \in V_1$

Example of graph G^*

$$V_1 = \{1, 2, 3, 4, 5\}$$

$$V_2 = \{[12345], [1234], [1245], [134], [234], [245], [13], [15], [34], [1], [3]\}$$



Proof of $\frac{3r}{2} \leq s$ using the graph G^*

V_1 LD-code and \nexists LD-code with vertices in both stable sets
 $\Rightarrow G^*$ satisfies:

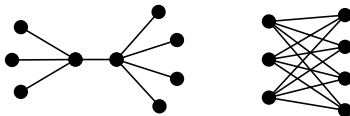
- $|V(G^*)| = s + 1$, $|E(G^*)| \geq 2r$
- G^* is bipartite
- incident edges have different labels
- walks contain an even number of edges with label v ,
 $\forall v \in V_1$, iff they are closed
- G^* contains a subgraph H^* of size $2r$ such that all its connected components are cactus
- G with cc. cactus with no C_3 , then $|E(G)| \leq \frac{4}{3}(|V(G)| - 1)$.

$$\Rightarrow s \geq \frac{3r}{2}$$

Bipartite graphs with $\lambda(\overline{G}) - \lambda(G) \in \{-1, 0, 1\}$

Given integers r, s , $3 \leq r \leq s$ there are bipartite graphs G with stable parts V_1, V_2 satisfying $|V_1| = r, |V_2| = s$ and:

- $\lambda(\overline{G}) = \lambda(G) - 1$: **double star** $K_2(r-1, s-1)$
- $\lambda(\overline{G}) = \lambda(G)$: **complete bipartite graphs** $K(r, s)$

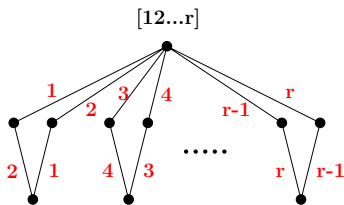


- $\lambda(\overline{G}) = \lambda(G) + 1$, if $\frac{3r}{2} + 1 \leq s \leq 2^r - 1$: **$G(r, s)$**

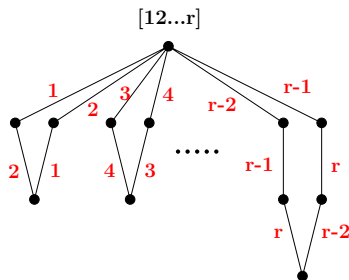
Bipartite graphs with $\lambda(\overline{G}) = \lambda(G) + 1$

- (r, s) , $r, s \in \mathbb{N}$, $3 \leq r$ and $\frac{3r}{2} + 1 \leq s \leq 2^r - 1$,
 $\exists G(r, s)$ bipartite graph such that $\lambda(\overline{G}) > \lambda(G)$.

$V_1 = \{1, 2, \dots, r\}$ and $s = \left\lceil \frac{3r}{2} + 1 \right\rceil$:



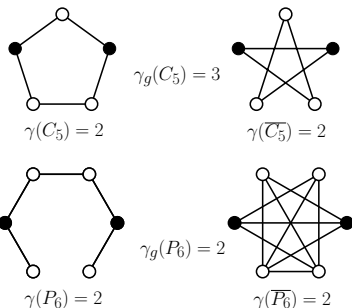
G^* , r even



G^* , r odd

Global domination [Brigham and Carrington, 1998]

- $S \subseteq V$, *global dominating set* of G :
 a dominating set of G and of \overline{G}
- *global domination number*, $\gamma_g(G)$:
 minimum size of a global dominating set of G



Global LD-number

- *global LD-number of G* :
minimum size of a global LD-set
- ✓ $\lambda_g(G) = \lambda_g(\overline{G})$

For any graph $G = (V, E)$,

- ▶ $\lambda(\overline{G}) \leq \lambda_g(G) \leq \lambda(G) + 1$
- ▶ $\lambda_g(G) = \lambda(G) \iff G$ has a global LD-code
 $\lambda_g(G) = \lambda(G) + 1 \iff$ every LD-code of G is non-global
- ▶ $\lambda(\overline{G}) \neq \lambda(G) \Rightarrow \lambda_g(G) = \max\{\lambda(G), \lambda(\overline{G})\}$
 $\lambda(\overline{G}) = \lambda(G) \Rightarrow \lambda_g(G) \in \{\lambda(G), \lambda(G) + 1\}$

Some values

	P_n ($n \geq 7$)	C_n ($n \geq 7$)	W_n ($n \geq 8$)	K_n ($n \geq 3$)
$\lambda(G)$	$\lceil \frac{2n}{5} \rceil$	$\lceil \frac{2n}{5} \rceil$	$\lceil \frac{2n-2}{5} \rceil$	$n-1$
$\lambda(\overline{G})$	$\lceil \frac{2n-2}{5} \rceil$	$\lceil \frac{2n-2}{5} \rceil$	$1 + \lceil \frac{2n-4}{5} \rceil$	n
$\lambda_g(G)$	$\lceil \frac{2n}{5} \rceil$	$\lceil \frac{2n}{5} \rceil$	$1 + \lceil \frac{2n-4}{5} \rceil$	n

	$K_{1,n-1}$ ($n \geq 4$)	$K_{r,s}$ ($2 \leq r \leq s$)	$K_2(r,s)$ ($2 \leq r \leq s$)
$\lambda(G)$	$n-1$	$n-2$	$n-2$
$\lambda(\overline{G})$	$n-1$	$n-2$	$n-3$
$\lambda_g(G)$	$n-1$	$n-2$	$n-2$